Induction Principle

# Natural Numbers and Induction

#### Ming Hsiung mingshone@163.com

School of Philosophy and Social Development South China Normal University

Set Theory, Spring 2022

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## The Set of Natural Numbers

# Induction Principle

**③** The Recursion Theorem

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## The Set of Natural Numbers

# **2** Induction Principle

**3** The Recursion Theorem

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## Recall

- Ø
- {Ø}
- $\{\emptyset, \{\emptyset\}\}$
- $\{\varnothing, \{\varnothing\}, \{\varnothing, \{\varnothing\}\}\}$
- .....

$$S(x) =_{\mathrm{df}} x \cup \{x\}$$

#### • 0

• 
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#### • Is there any infinite set?

- Answer: the Axiom of Infinity is indpendent of the other axioms of ZFC.
- The Axiom of Infinity: There is a set I such that  $0 \in I$  and if  $x \in I$ , then  $S(x) \in I$ .
- Definition: a set I is called **inductive**, if  $0 \in I$  and if  $x \in I$ , then  $S(x) \in I$ .

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## The set $\mathbb{N}$

- Definition: ℕ =<sub>df</sub> {*x*|*x* belongs to every inductive set}. The elements of ℕ are called the natural numbers.
- Formally,  $\mathbb{N} =_{df} \dots$
- Fact: N itself is an inductive set, and it is a subset of every inductive set.

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## The Set of Natural Numbers

# 2 Induction Principle

**3** The Recursion Theorem

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## **Induction Principle**

- The Induction Principle: Let P(x) be a property (possibly with parameters). Assume that
  - (a) P(0) holds.
  - (b) For all  $n \in \mathbb{N}$ , P(n) implies P(S(n)).

Then P(n) holds for any  $n \in \mathbb{N}$ .

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- Definition: The relation < on  $\mathbb{N}$  is defined by: m < n iff  $m \in n$ .
- Definition:  $m \leq n =_{df} m \in n \lor m = n$ .
- Fact: For any  $n \in \mathbb{N}$ ,  $0 \leq n$ .

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- Definition: The relation < on  $\mathbb{N}$  is defined by: m < n iff  $m \in n$ .
- Definition:  $m \leq n =_{df} m \in n \lor m = n$ .
- Fact: For any  $n \in \mathbb{N}$ ,  $0 \leq n$ .

#### • Theorem: < is a linear ordering on $\mathbb{N}$ .

#### **Induction Principle, Second Version**

• The Induction Principle, Second Version: Let P(x) be a property (possibly with parameters). Assume that for  $n \in \mathbb{N}$ ,

If P(k) holds for all k < n, then P(n).

Then P(n) holds for any  $n \in \mathbb{N}$ .

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- Definition: A relation R on A is well-founded, if for every nonempty subset B of A, there exists a R-minimal element of B, i.e., an element x ∈ B such that yRx fails for any y ∈ B.
- Fact: A relation R on A is well-founded, iff no elements x<sub>0</sub>, x<sub>1</sub>, ..., of A can form a decreasing chain, that is, x<sub>1</sub>Rx<sub>0</sub>, x<sub>2</sub>Rx<sub>1</sub>, .... Proof. ⇒:



Note: The  $\Leftarrow$  side needs to employ the Axiom of Choice (AC). People usually use AC without realizing it!

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#### • Theorem: < is a well-founded relation on $\mathbb{N}$ .



## The Set of Natural Numbers

# **2** Induction Principle

# The Recursion Theorem

#### Informally,

- $a_0, a_1, \ldots, a_{n-1}$  is a sequence of length n.
- $a_0, a_1, \ldots, a_n, \ldots$  is an infinite sequence (of length  $\omega$ ).

#### Formally,

- A finite sequence of elements of A is a function from n to A for some n. n is called the length of this sequence. If n = 0, we get the empty sequence.
- An (countably) **infinite sequence of elements of** *A* is a function from ℕ to *A*.

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- $1, 1, 2, 6, 24, \ldots, 1 \cdot 2 \cdot \ldots \cdot n, \ldots$
- Explicit Definition: *f* is a function such that f(n) = n!(stipulate: 0! = 1).
- Implicit Definition: f is a function such that f(0) = 1 and  $f(S(n)) = f(n) \cdot S(n)$ .<sup>1</sup>

<sup>1</sup>This is an example for illustrating the idea of recursion, as we still have not a well-defined notion of multiplication for natural numbers.  $\langle \Box \rangle \langle \Box \rangle \langle \Xi \rangle \langle \Xi \rangle \langle \Xi \rangle \langle \Xi \rangle$ 

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## **Recursion on Natural Numbers**

Theorem (The Recursion Theorem)

For any set A, any a ∈ A, and any function g : A × N → A, there exists a unique function f : N → A such that
(a) f(0) = a;

(b) f(S(n)) = g(f(n), n) for all  $n \in \mathbb{N}$ .

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Application 1: For every m ∈ N, there exists a unique function f<sub>m</sub> : N → N such that
(a) f<sub>m</sub>(0) = m;

(b) 
$$f_m(S(n)) = S(f_m(n))$$
 for all  $n \in \mathbb{N}$ .

Proof.

#### • Definition.

$$m+n =_{\mathrm{df}} f_m(n).$$

• Thus:

(a) 
$$m + 0 = m$$
;  
(b)  $m + S(n) = S(m + n)$  for all  $m, n \in \mathbb{N}$ .

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• Prove: 
$$m + 1 = S(m)$$

From now on, we can use m + 1 instead of S(m).

#### Theorem (The Recursion Theorem with parameters)

For any  $a : P \to A$  and any function  $g : P \times A \times \mathbb{N} \to A$ , there exists a unique function  $f : P \times \mathbb{N} \to A$  such that

(a) 
$$f(p,0) = a(p);$$

(b) f(p, S(n)) = g(p, f(n), n) for all  $n \in \mathbb{N}$  and  $p \in P$ .

Application 2: There exists a unique function *f* : N × N → N such that

(a) f(m, 0) = 0;(b) f(m, S(n)) = f(m, n) + m for all  $m, n \in \mathbb{N}.$ 

Proof.

#### • Definition.

$$m \cdot n =_{\mathrm{df}} f(m, n).$$

• Thus:

(a) 
$$m \cdot 0 = m$$
;  
(b)  $m \cdot S(n) = m \cdot n + m$  for all  $m, n \in \mathbb{N}$ .

• Application 3: There exists a unique function  $f: \mathbb{N} \to \mathbb{N}$  such that

(a) 
$$f(0) = 1$$
;  
(b)  $f(S(n)) = f(n) \cdot n$  for all  $n \in \mathbb{N}$ .

Proof.

• Definition.

$$n! =_{\mathrm{df}} f(n).$$

• Thus:

(a) 
$$0! = 1;$$
  
(b)  $(n+1)! = n! \cdot n$  for all  $n \in \mathbb{N}$ 

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#### Theorem (The Recursion Theorem)

For any set A, any  $a \in A$ , and any function  $g : A \times N \to A$ , there exists a unique function  $f : \mathbb{N} \to A$  such that

(a) f(0) = a;

(b) f(S(n)) = g(f(n), n) for all  $n \in \mathbb{N}$ .

Proof.

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## Homework

- How to define exponentiation of natural numbers? Justify your definition.
- How to define (by recursion) Fibonacci function f : N → N, which meets the condition:

$$\begin{cases} f(0) = 1\\ f(1) = 1\\ f(n+2) = f(n+1) + f(n), n \ge 0 \end{cases}$$

Hint: Refer to Hrbaeck & Jech (3 ed.), pp. 50.

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# Thanks for your attention! Q & A

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